IDUG

Unleashing the Potential of Columnar Tables in Db2

Satya Krishnaswamy, IBM

Central Canada Db2 Users Group

Toronto

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AGENDA

Columnar Data Engine (CDE) in Db2

Storage and compression of Columnar Table

Recent improvements

What's coming

Creating a Column-Organized Table

CREATE TABLE sales_col (c1 INTEGER NOT NULL, c2 INTEGER, . . . PRIMARY KEY (c1)) ORGANIZE BY COLUMN;

• If dft_table_org = COLUMN

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- ORGANIZE BY COLUMN is the default and can be omitted
- Use ORGANIZE BY ROW to create row-organized tables
- DB2_WORKLOAD=ANALYTICS Sets Everything You Need

What you see in the Db2 catalog: TABLEORG

- Which tables are column-organized?
 - Column In syscat.tables: TABLEORG



Columnar storage in Db2 (conceptual)

Age	Address	City	State	Zip Code
47	18 Main Street	Springfield	MA	01111
32	455 N. 1 st St.	San Jose	СА	95113
55	911 Elm St.	Toledo	ОН	43601
22	300 Grand Ave	Los Angeles	СА	90047
43	404 Escuela St.	Los Angeles	CA	90033
29	166 Elk Road #47	Beaverton	OR	97075
78	5661 Bloom St.	Raleigh	NC	27605
35	8883 Loghorn Dr.	Tucson	AZ	85701
80	14 California Blvd.	Pasadena	CA	91117
61	1114 Apple Lane	Cupertino	СА	95014
	Age 47 32 55 22 43 29 78 35 80 61	Age Address 47 18 Main Street 32 455 N. 1 st St. 55 911 Elm St. 55 911 Elm St. 22 300 Grand Ave 43 404 Escuela St. 29 166 Elk Road #47 78 5661 Bloom St. 35 8883 Loghorn Dr. 80 14 California Blvd. 61 1114 Apple Lane	AgeAddressCity4718 Main StreetSpringfield32455 N. 1 st St.San Jose55911 Elm St.Toledo22300 Grand AveLos Angeles43404 Escuela St.Los Angeles29166 Elk Road #47Beaverton785661 Bloom St.Raleigh358883 Loghorn Dr.Tucson8014 California Blvd.Pasadena611114 Apple LaneCupertino	AgeAddressCityState4718 Main StreetSpringfieldMA32455 N. 1 st St.San JoseCA55911 Elm St.ToledoOH22300 Grand AveLos AngelesCA43404 Escuela St.Los AngelesCA29166 Elk Road #47BeavertonOR785661 Bloom St.RaleighNC358883 Loghorn Dr.TucsonAZ8014 California Blvd.PasadenaCA611114 Apple LaneCupertinoCA

Row-Organized Table Format

• Traditional approach: data stored in row format

John Piconne	47	18 Main Street	Springfield	MA	01111	3
Susan Nakagawa	32	32 455 N. 1 st St. San Jose		CA	95113	Page
Sam Gerstner	55	911 Elm St.	Toledo	ОН	43601	
Chou Zhang	22	300 Grand Ave	Los Angeles	СА	90047	1. Alternational Action of the second
Mike Hernandez	43	404 Escuela St.	Los Angeles	CA	90033	Page
Pamela Funk	29	166 Elk Road #47	Beaverton	OR	97075	
Rick Washington	78	5661 Bloom St.	Raleigh	NC	27605	
Ernesto Fry	35	8883 Longhorn Dr.	Tucson	AZ	85701	Page
Whitney Samuels	80	14 California Blvd.	Pasadena	CA	91117	гауе
Carol Whitehead	61	1114 Apple Lane	Cupertino	CA	95014	
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• Each page contains 1 or multiple rows (all columns)

Tuple Sequence Number (TSN)

		Column	(Colum	n	Column		Column	(Column	Column	
	TSN	Group 1	_ (Group	2	Group 3		Group 4	(Group 5	Group 6	11111114
	0	John Piconne		47		18 Main Street		Springfield		MA	01111	tun la
	1	Susan Nakagawa		32		455 N. 1 st St.		San Jose		CA	95113	
TSN =	2	Sam Gerstner		55		911 Elm St.		Toledo		ОН	43601	none
Tunle	3	Chou Zhang		22	30	300 Grand Ave		Los Angeles		CA	90047	< page
Sequence	4	Mike Hernandez		43		404 Escuela St.		Los Angeles		CA	90033	
Number	5	Pamela Funk		29		166 Elk Road #47		Beaverton		OR	97075	
	6	Rick Washington		78		5661 Bloom St.		Raleigh	Ł	NC	27605	
	7	Ernesto Fry		35		8883 Longhorn Dr.		Tucson		AZ	85701	
	8	Whitney Samuels	80			14 California Blvd.	1	Pasadena	1	CA	91117	ſ
	9	Carol Whitehead	L	61		1114 Apple Lane		Cupertino		CA	95014	
	10		À	-011		NUMBER AND	Т		ľ			<
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- Each tuple (row) in the table is assigned a TSN, which is similar to a Row ID. •
- A tuplet is the projection of a column group on a tuple
- TSNs are used to stitch tuplets together during query processing •
- The "start TSN" of each page is stored in the Page Map Index (PMI) for lookup of values • by TSN

Extents and Pages

- Each column group has their own data pages to store column data
- Data pages are grouped into extents
- Extents, rather than individual pages, are the basic unit of data to be allocated and freed when required
- Typically, four pages will make up one extent



Column Groups and NULL values

- Each column in the table belongs to exactly one column group
- Currently, the only reason for a column group to hold multiple columns is if a column is nullable.
- The column in the table will be represented in the column group as a nullable column holding the data and an internal null indicator column.
 - The null indicator column contains 1-bit values indicating if the value for the corresponding column in the table is null.
- Null indicator values are either packed together with the data values in a single bank or stored in a separate bank within a region.

Null values example

Table (External Representation):

Name (NOT NULL)	Age	Address	City	State	Zip Code
John Piconne	NULL	18 Main Street	NULL	MA	01111
Susan Nakagawa	32	NULL	San Jose	NULL	95113

Table (Internal Representation):

Name	Age	NI (Age)	Address	NI (Address)	City	NI (City)	State	NI (State)	Zip Code	NI (Zip Code)
John Piconne	?	1	18 Main Street	0	?	1	MA	0	01111	0
Susan Nakagawa	32	0	?	1	San Jose	0	?	1	95113	0

? = Value that can be encoded to the least amount of bits

Synopsis Table

 Meta-data that describes which ranges of values exist in which parts of the user table

SYN130330165216275152 SALES COL

	ISMMAA	S_DATEMIN	S_DATEMAX	• • •
0	1023	2005-03-01	2006-10-17	
1024	2047	2006-08-25	2007-09-15	• • •

User table: **SALES**_COL

	S_DATE	QTY	•••
	2005-03-01	176	•••
	2005-03-02	85	• • •
	2005-03-02	267	
	2005-03-04	231	
111			
	3 • • •		

TSN = Tuple Sequence Number

Enables Db2 to skip portions of table when scanning data during query

1023 1024

2047

- Benefits from data clustering, loading pre-sorted data
- Predicate WHERE S_DATE = 2007-01-01 would skip first range
- Predicate WHERE S_DATE = 2006-09-12 would scan both ranges



Synopsis Table: Example

Base Table:

TSN	Name	Age
0	John Piconne	47
1	Susan Nakagawa	32
2	Sam Gerstner	55
3	Chou Zhang	22
4	Mike Hernandez	43
5	Pamela Funk	29
6	Rick Washington	78
7	Ernesto Fry	35
8	Whitney Samuels	80
9	Carol Whitehead	61

Synopsis Table (Synopsis tuple covers 5 base table tuples):

TSN MIN	TSN MAX	NAMEMIN	NAMEMAX	AGE MIN	AGE MAX
0	4	Chou Zhang	Susan Nakagawa	22	55
5	9	Carol Whitehead	Whitney Samuels	29	80

What you see in the Db2 catalog: Synopsis Tables

- Each columnar table has a corresponding synopsis table with a few exceptions
- Automatically created/maintained
- Queries use the synopsis table to determine if it can skip ranges of rows when evaluating predicates

When Data is Encoded vs. Unencoded

• Most columns have a mix of both encoded and unencoded values.

Encoded	Unencoded
 A column value is stored encoded if a dictionary entry can be used to encode that data Dictionary entry can be used Internal columns like NULL indicator columns, which are encoded using 1 bit DECIMAL columns, minus coding if precision <=18 (imposed by 64-bit code size limit) 	 A value is stored unencoded if it cannot be represented in the dictionary Values that occur infrequently (e.g., value occurs once, inefficient to create a dictionary value) Values that were pruned from the histogram due to memory constraints Values that were INSERTed/LOADed after dictionary creation, not covered by existing dictionary (static)

Encoding differs from **compression**, a value can be unencoded and still be compressed with ..., but the value needs to be decompressed before queries can be performed on it. This is why we refer to pure dictionary encoded and encoding types as *actionable* compression.

Columnar Compression

• Columnar compression dictionaries are used to compress data in a column of a columnar table by mapping repeated byte patterns to much smaller symbols, which then replace the longer byte patterns in the table.

Frequency compression

- Most common values use fewest bits
- Encoding Schemes
 - Pure dictionary coding
 - Prefix coding
 - Minus coding

Since all data in a column are of the same data type, we can **improve the compression ratio** by exploiting the skew in data distribution.

Column compression differs from row compression in that we **map values to dictionary codes**.

Benefits

- **Save disk space** by reducing total amount of **buffer pool** pages needed to store data.
- **Faster queries**, a compressed table need fewer I/O operations to access the same amount of data.

Column-level Dictionaries are Static

• Once created, column-level dictionaries are never updated

- REORG does not rebuild column-level dictionaries
- Row organized tables can use REORG to rebuild the dictionary
- The user must unLOAD and reLOAD the table to rebuild the dictionary

Page compression reduces the need to rebuild column dictionaries

- New values not covered by the initial column-level dictionary can still be compressed by the page-level dictionary
- This reduces deteriorating the compression ratio over time
- When a page fills up, decide whether to do page compression



Best Practices – Enable All Optional BLU Storage Enhancements

Feature	Description	Release	Db2	Db2 Warehouse
Page-based String Compression Type 1	Improves compression for high cardinality string columns with repeating portions of strings that are not encoded by existing compression algorithms	11.5.4	Need Registry Variable to Enable Registry Variable (2)	Enabled by Default
Page-Based String Compression Type 2	Improved compression when strings (within a page) contain 16 or less unique characters, works well for hex, numeric items like phone numbers, dates, dollar values when stored as strings	11.5.4	Need Registry Variable to Enable Registry Variable (2)	Enabled by Default
Deferred Synopsis Tuple Creation for Small Base Tables	Reduces synopsis table storage consumption for small base tables	11.5.4	Need Registry Variable to Enable Registry Variable (3)	Enabled by Default

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Best Practices – Enable All Optional BLU Storage Enhancements

Feature	Description	Release	Db2	Db2 Warehouse
Reorg Table Recompress Enhancement	Improves performance of Reorg Table Recompress and applies page-based string compression during Reorg Table Recompress	11.5.5	Enabled by Default	Enabled by Default
Trickle Feed Insert Enhancement	Speeds up trickle feed insertion, reduces the memory footprint and size of small tables.	11.5.6	Not available	Enabled by Default
		11.5.7	Need Registry Variable to Enable Registry Variable (1)	Enabled by Default
LOAD Utility Enhancement	Improves overall LOAD processing, and also if the previously mentioned string compression algorithms are enabled allows LOAD to use them	11.5.8	Enabled by Default	Enabled by Default

7 Big Ideas

- Simple to Implement and Use (Load & Go)
- Data Remains Compressed During Evaluation
- Multiply the Power of the CPU
- Core-Friendly Parallelism
- Work performed directly on columns
- Scan-Friendly Memory Caching
- Data skipping

7 Big Ideas: How Db2 with BLU Acceleration Helps ~Sub second 10TB query – An Optimistic Illustration

- The system 32 cores, 10TB table with 100 columns, 10 years of data
- The query: SELECT COUNT (*) from MYTABLE where YEAR = '2010'
- The optimistic result: sub second 10TB query! Each CPU core examines the equivalent of just 8MB of data



Thank You

Speaker: Company: Email Address: Satya Krishnaswamy IBM satya.ksr@us.ibm.com